

# Exercise Sheet 3

Solutions due: 23.06.2016

Rev. 12709

## Exercise 1 Frame Definability

1. Recall that we refer to the normal modal logic generated by the Löb-axiom

$$(L) \quad \Box(\Box p \rightarrow p) \rightarrow \Box p$$

as *provability logic* and that we denote this logic by  $\mathbf{L}$ . We know that  $\mathbf{L}$  is sound and complete w.r.t. the class  $FTT$  of all frames that have a finite transitive tree as relation.

Prove Lemma 3.37 from the lecture, i.e. show that there is no FOL-formula that locally corresponds to (L).

2. Proposition 3.39 from the lecture implies that modal formulas locally correspond to their respective *second-order* translation. Prove Proposition 3.39 by showing that for all modal formulas  $\varphi$  that use  $n$  propositional variables  $p_1, \dots, p_n$ , all frames  $\mathcal{F} = (\mathcal{W}, \mathcal{R})$  and all worlds  $w \in \mathcal{W}$ ,

$$(a) \quad \mathcal{F}, w \Vdash \varphi \text{ iff } \mathcal{F}, \eta \models_{MSO} \forall P_1, \dots, P_n. ST_x(\varphi) \text{ for } \eta = [w/x],$$

$$(b) \quad \mathcal{F} \Vdash \varphi \quad \text{iff } \mathcal{F} \models_{MSO} \forall P_1, \dots, P_n. ST_x(\varphi).$$

## Exercise 2 Filtration

Prove Item 2. of Theorem 4.2 from the lecture, i.e. prove that for all finite sets of formulas  $\Gamma$  that are closed under taking subformulas, all models  $\mathcal{M}$ , all formulas  $\varphi \in \Gamma$  and all worlds  $w$  of  $\mathcal{M}$ , we have

$$\mathcal{M}, w \models \varphi \text{ iff } \mathcal{M}^f, [w] \models \varphi,$$

where  $\mathcal{M}^f$  denotes the filtration of  $\mathcal{M}$  through  $\Gamma$ .

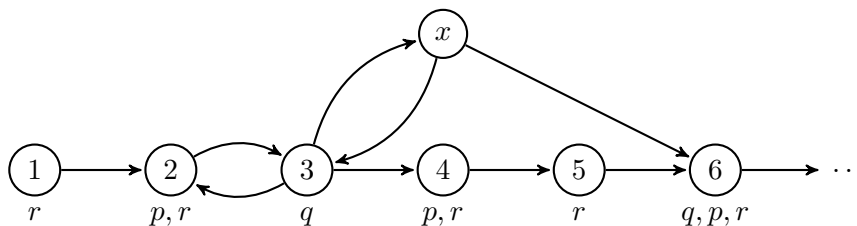
## Exercise 3 Selection

1. Finish the proof of Theorem 4.6 from the lecture by proving that for all formulas  $\psi$ ,

$$\mathcal{M}, w \models \psi \text{ iff } \mathcal{M}^{\psi, w}, w \models \psi,$$

where  $\mathcal{M}^{\psi, w}$  is obtained from  $\mathcal{M}$  by selection, i.e. by restricting the model to the worlds that are relevant for the satisfaction of  $\psi$  at world  $w$ .

2. Consider the model  $\mathcal{M} = (\mathcal{W}, \mathcal{R}, \mathcal{V})$



given by

$$\begin{aligned} \mathcal{W} &= \{x\} \cup \{i \mid i > 0\} \\ \mathcal{R} &= \{(i, i+1) \mid i > 0\} \cup \{(3, 2), (3, x), (x, 3), (x, 6)\} \\ \mathcal{V}(p) &= \{j \mid j \text{ is even}\} \\ \mathcal{V}(q) &= \{j \mid \exists i \in \mathbb{N}. j = 3i\} \\ \mathcal{V}(r) &= \mathbb{N}_{>0} \setminus \{3\} \end{aligned}$$

Let  $\psi = \diamond(\neg\diamond r \vee \diamond(\diamond p \wedge \diamond q))$  and compute the result of applying the selection function from the proof of Theorem 4.6 to  $\psi$  and 1, i.e. compute  $s(\psi, 1)$ . Use this set to define the model  $\mathcal{M}^{\psi, 1}$  that is obtained from  $\mathcal{M}$  by selection.

## Exercise 4 Model size

Theorem 4.12 from the lecture states that there is for each  $\Lambda \in \{\mathbf{K}, \mathbf{T}, \mathbf{B}, \mathbf{K4}, \mathbf{S4}\}$  and each  $n \in \mathbb{N}$  a formula  $\varphi_n$  s.t.

1.  $|\varphi_n| \in \mathcal{O}(n^2)$ ,
2.  $\varphi_n$  is  $\Lambda$ -satisfiable, and
3. for all models  $\mathcal{M} = (\mathcal{W}, \mathcal{R}, \mathcal{V})$  of  $\varphi_n$ , we have  $|\mathcal{W}| \geq 2^n$ .

Prove this by devising a series of formulas  $\varphi_n$  such that for each  $n \in \mathbb{N}$ ,  $\varphi_n$  defines a full binary tree of depth  $n$  and by verifying properties 1., 2. and 3. for your series of formulas.

## Exercise 5 Sahlqvist Formulas

5 points

1. Are the following formulas Sahlqvist formulas?

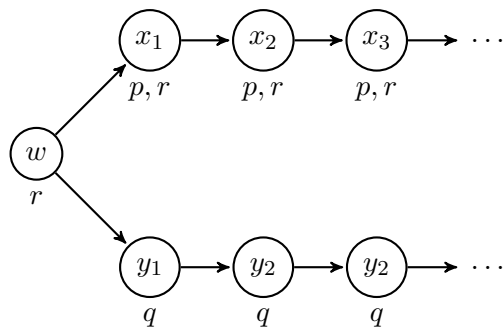
$$\begin{aligned} \psi_1 &:= \diamond \Box p \rightarrow \Box \diamond p \\ \psi_2 &:= \Box \diamond p \rightarrow \diamond \Box p \\ \psi_3 &:= (p \wedge \diamond \diamond p) \rightarrow \diamond p \\ \psi_4 &:= (p \wedge \Box p \wedge \Box \Box p) \rightarrow \diamond p \end{aligned}$$

2. For any formula from the previous item that is a Sahlqvist formula, define a first-order formula that expresses the same property on frames.

## Exercise 6 Filtration

9 points

1. Consider the model  $\mathcal{M} = (\mathcal{W}, \mathcal{R}, \mathcal{V})$



given by

$$\begin{aligned}\mathcal{W} &= \{w\} \cup \{x_i \mid i > 0\} \cup \{y_i \mid i > 0\} \\ \mathcal{R} &= \{(w, x_1), (w, y_1)\} \cup \{(x_i, x_{i+1}) \mid i > 0\} \cup \{(y_i, y_{i+1}) \mid i > 0\} \\ \mathcal{V}(p) &= \{x_i \mid i > 0\} \\ \mathcal{V}(q) &= \{y_i \mid i > 0\} \\ \mathcal{V}(r) &= \{w\} \cup \{x_i \mid i > 0\}\end{aligned}$$

Let  $\Gamma = \{\diamond p, \diamond q, \diamond r\}$ . Compute  $\sim_\Gamma$  and use this relation to find

- the smallest filtration of  $\mathcal{M}$  through  $\Gamma$ , and
- the largest filtration of  $\mathcal{M}$  through  $\Gamma$ .

Is there an upper bound on the number of worlds that filtrations of an arbitrary model  $\mathcal{M}'$  through  $\Gamma$  may have?

2. Prove that for all models  $\mathcal{M}$  that are based on a symmetric frame and all sets of formulas  $\Gamma$ ,
  - (a) the largest filtration of  $\mathcal{M}$  through  $\Gamma$  *not* necessarily has a symmetric relation.
  - (b) the smallest filtration of  $\mathcal{M}$  through  $\Gamma$  has a symmetric relation.

To prove Item (a), it suffices to find a counterexample. Use Item (b) to adapt the proof of Theorem 4.5 from the lecture to show that  $SAT(\mathbf{B})$  and  $VAL(\mathbf{B})$  are decidable.

## Exercise 7 *Tableaux*

**6 points**

Use the Tableau-algorithm from the lecture to decide whether the following formulas are *valid*:

1.  $\Box(p \rightarrow (\diamond q \wedge \diamond r)) \wedge \diamond p \wedge \diamond q \wedge \Box\Box((q \wedge r) \rightarrow p)$
2.  $\Box(p \rightarrow (\diamond q \vee \diamond r)) \wedge (\diamond p \vee \diamond \neg p) \wedge \Box\Box(p \wedge \neg r)$