

Dual Bisimilarity Games

ÜSAME CENGİZ

The Setting

Given a coalgebra (X, γ) , are two states in X bisimilar?

States are bisimilar iff there exists a bisimulation...

What is a bisimulation? Huh?

Bisimulation

A bisimulation is an eq. rel R , s.t.

$$2 \times X^A$$

$$A \rightarrow PX$$

$$\begin{array}{ccc} x & R & y \\ a \downarrow & & \downarrow a \\ x' & R & y' \\ & \& & \end{array}$$

$$x \text{ final} \Leftrightarrow y \text{ final}$$

$$\begin{array}{ccc} x & R & y \\ a \downarrow & & \downarrow a \\ S & \tilde{R} & T \\ & \circ & \circ \end{array}$$

"Just lift R along functor"

$$\begin{array}{l} \forall x' \in S. \\ \exists y' \in T. x' R y' \\ \text{and } \forall y' \in T \\ \exists x' \in S. x' R y' \end{array}$$

The Standard Game

We play on an LTS $(A \rightarrow PX)$

Positions are pairs $(x, y) \in X * X$

I win
infinite games!

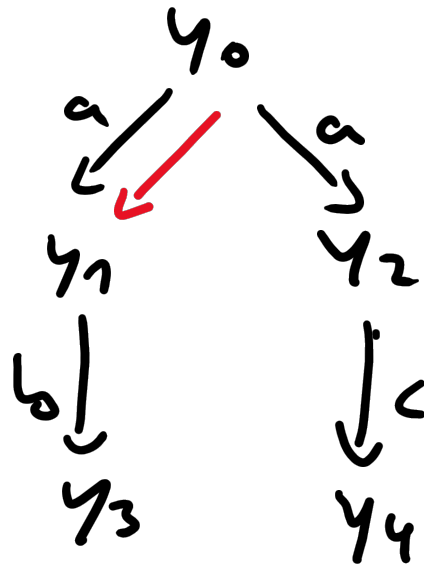
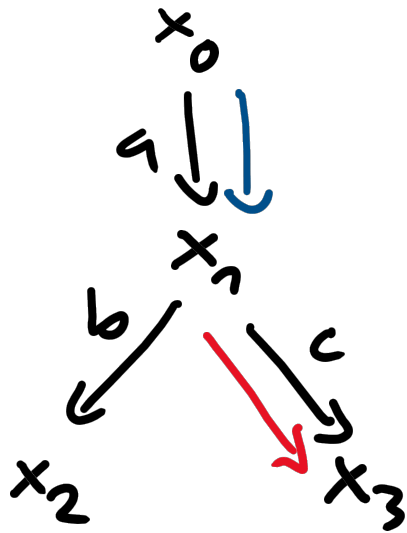
Players are Spoiler and Duplicator

$x \neq y!$

$x = y!$

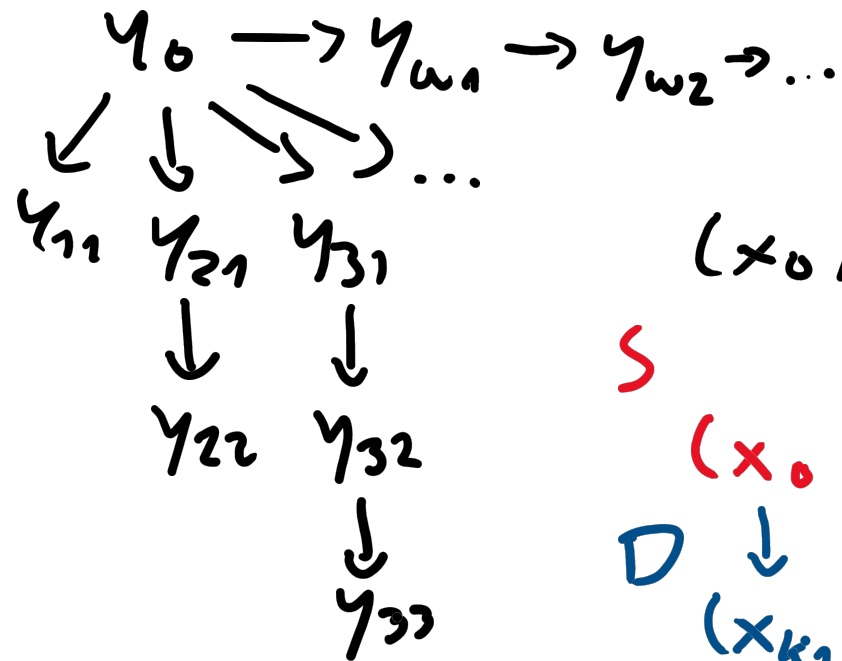
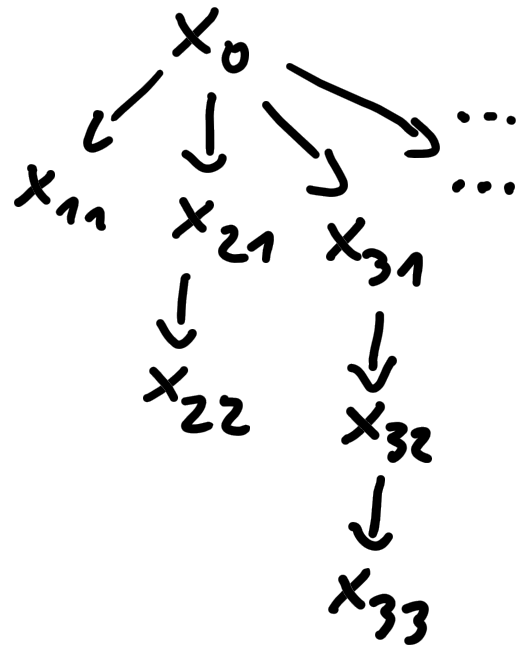
1. Spoiler chooses an $\overset{A}{\downarrow}$ -successor of either x or y
 2. Duplicator must copy said move, repeat
- Whoever cannot move loses!

The Standard Examples



(x_0, y_0)
 S $\downarrow a$
 (x_0, y_1)
 D $a \downarrow$
 (x_1, y_1)
 S $c \downarrow$
 (x_3, y_1)
 D $\neg \neg (u) \neg$

The Standard Examples



(x_0, y_0)
S $\downarrow$
 (x_0, y_{w1})
D $\downarrow$
 (x_{k1}, y_{w1})
S wins in k turns

The Quest for Generalization

Graded Semantics

Codensity Liftings

induce
games

(Ford et al.)
(Forster et al.)

(Komorida et al.)

„They seem almost dual to each other“

- Jonas Forster

(also everyone else I've talked to)

Games via Graded Semantics*

Position	PL	Move
$(x, y) \in X \times X$	D	"Equations" $R \subseteq X \times X$ s.t. $\gamma(x) =_R \gamma(y)$
$R \subseteq X \times X$	S	$(x', y') \in R$

$$\models X = P(A \times X)$$

local bisimulation
 δ^*

believable for
one step

Spoiler doesn't agree with
this equation

*simplified

Games via Graded Semantics*

For Trace Equivalence!

Position	PL	Move
$(x, y) \in PX \times PX$	D	"Equations" $R \subseteq PX \times PX$ s.t. $\gamma^{\#}(x) =_R \gamma^{\#}(y)$
$R \subseteq X \times X$	S	$(x', y') \in R$

$$\boxed{\models X = P(A \times X)}$$

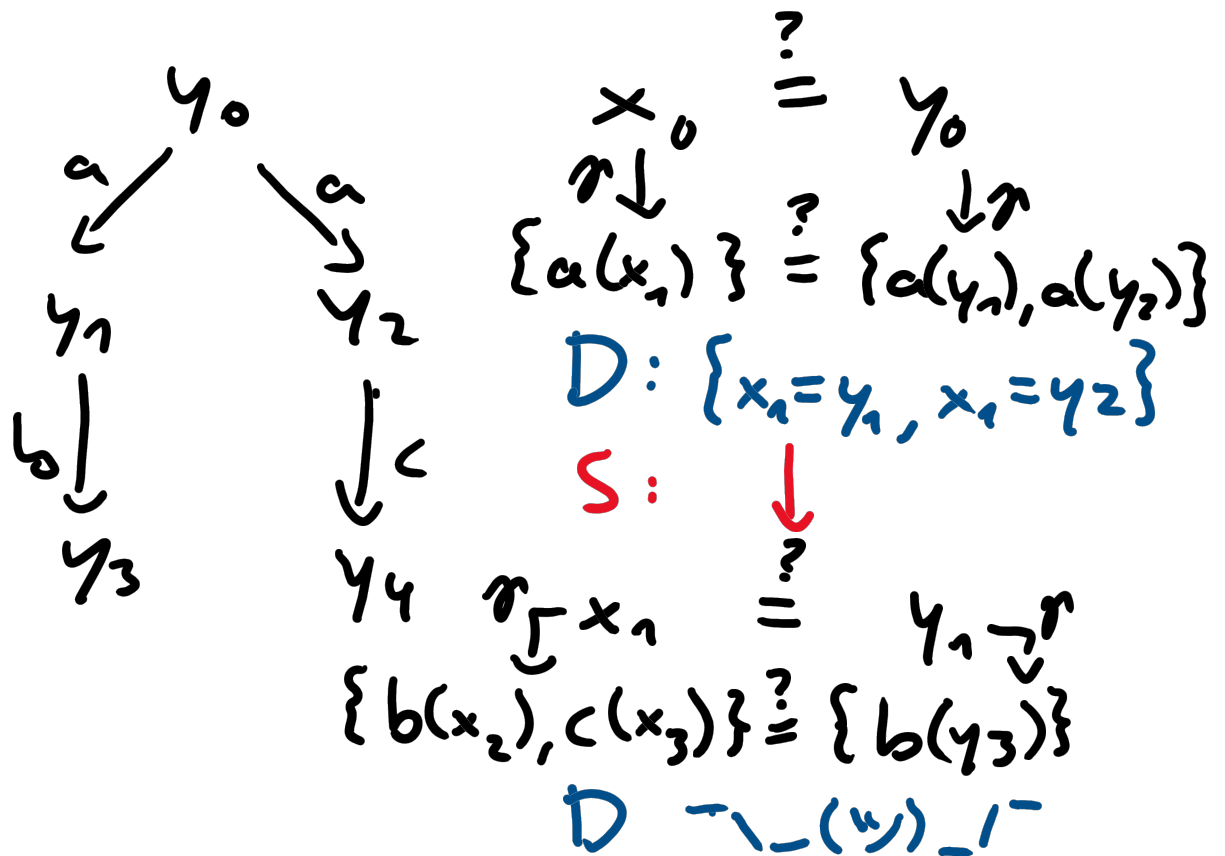
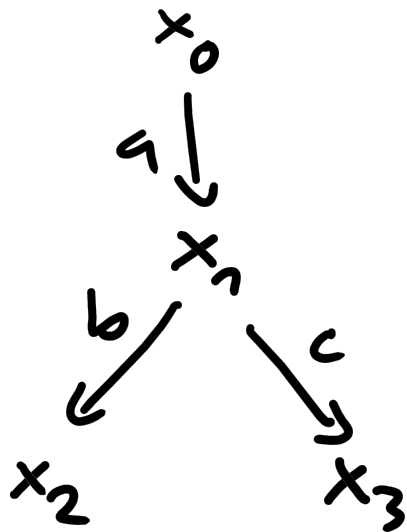
Just play on the
determinization

via
abstract
nonsense!

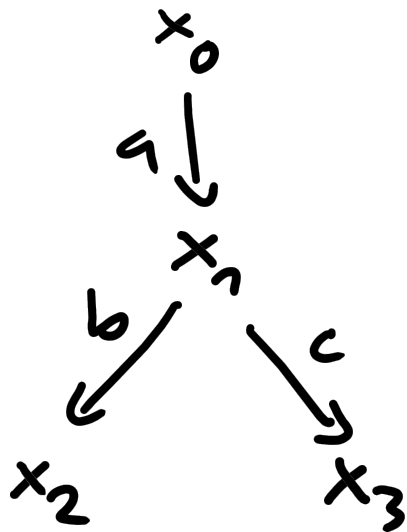
$x \sim_{tr} y$ iff D wins $(\{x\}, \{y\})$
 finite! $\nearrow$

*simplified

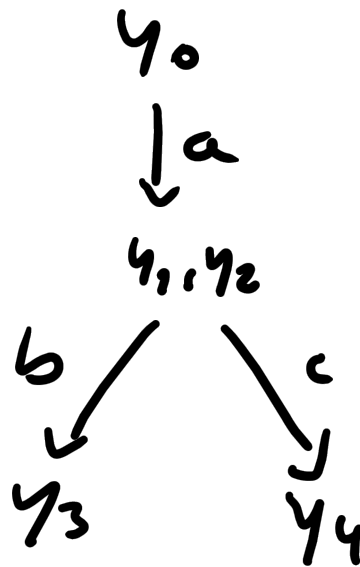
Playing the Game



Playing the Game For Trace Equivalence!



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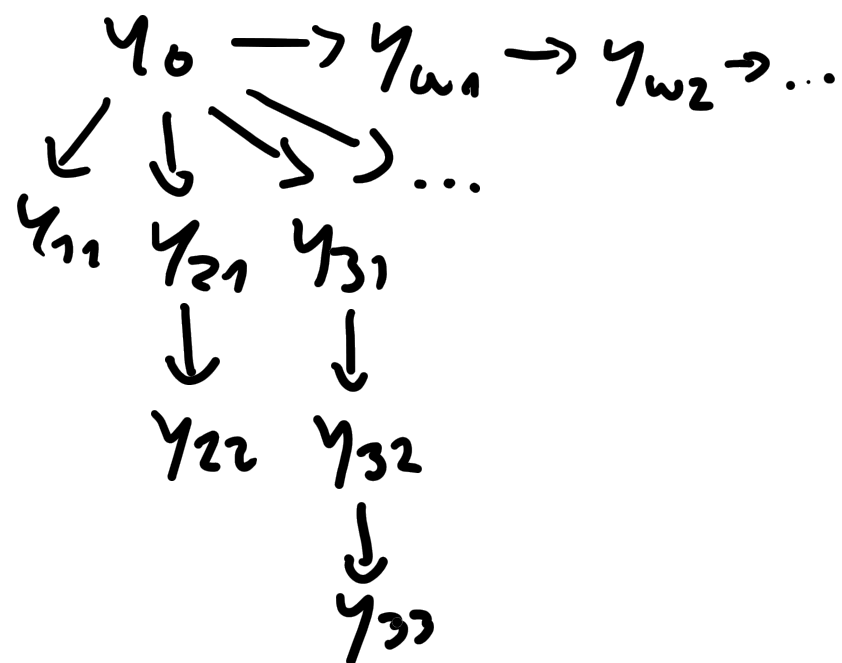
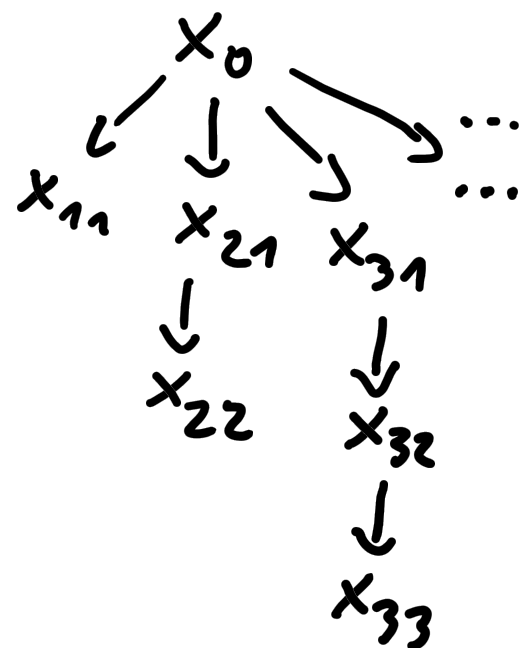


S can't choose $\leftarrow$

$$\begin{aligned}
 &\{x_0\} \stackrel{?}{=} \{y_0\} \\
 &\quad \downarrow \tau^\# \quad \quad \quad \downarrow \tau^\# \\
 &\{a(\{x_1\})\} \stackrel{?}{=} \{a(\{y_1, y_2\})\} \\
 &D: \{x_1\} = \{y_1, y_2\} \\
 &S: \quad \quad \quad \downarrow \\
 &\tau^\# \{x_2\} \stackrel{?}{=} \{y_3\} \tau^\# \\
 &\quad \quad \quad \downarrow \tau^\# \\
 &\quad \quad \quad \{ \} \stackrel{?}{=} \{ \}
 \end{aligned}$$

D:

Try both versions with this one



Games via Codensity Liftings*

Position	PL	Move
$(x, y) \in X \times X$	S	„Predicate“ $k : X \rightarrow 2$ s.t. $\leftarrow$ $\exists a. T \in Fk \circ \rho(x)(a)$ $\Downarrow$ $T \notin Fk \circ \rho(y)(a)$
$k : X \rightarrow 2$	D	(x', y') s.t. $k(x') \neq k(y')$

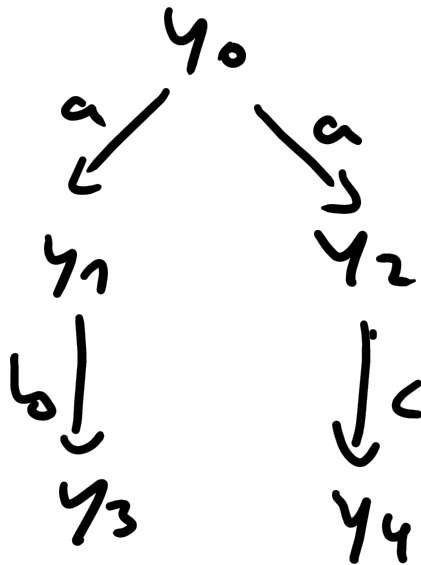
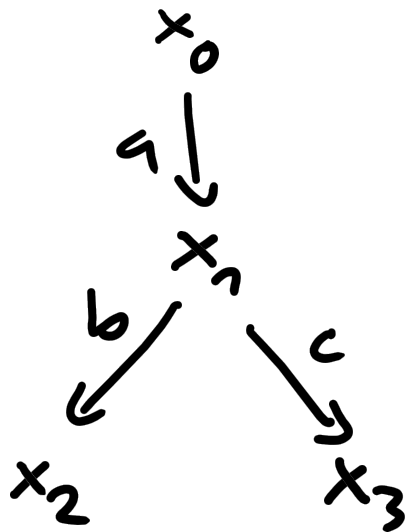
Duplicator thinks this isn't
respecting bisimilarity $\nearrow$

$$\boxed{\exists X = D(A \times X)}$$

Spoiler plays
a set of
distinguishing
behaviours

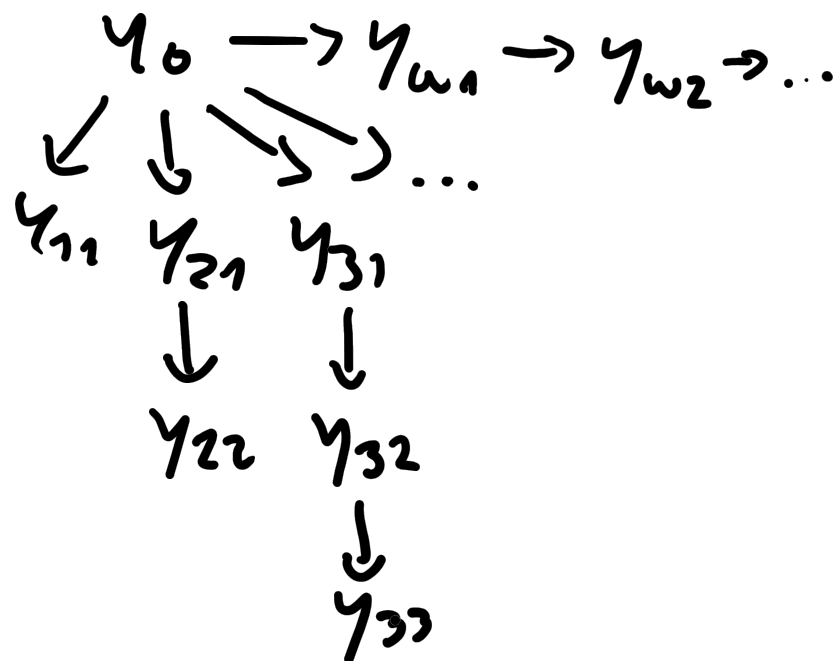
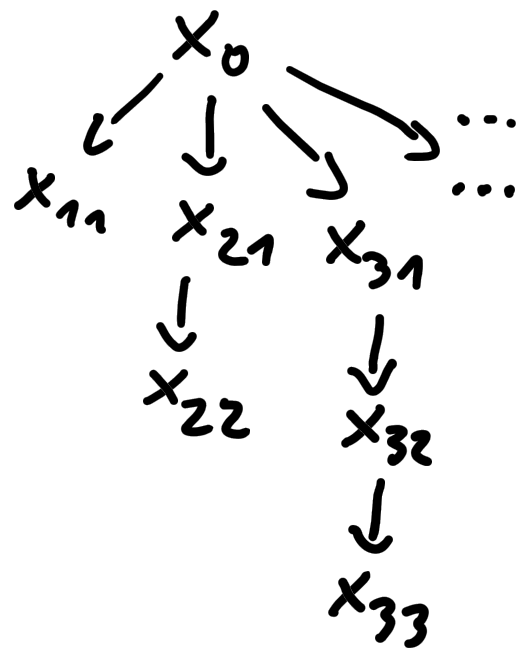
*simplified

Playing the Game



$x_0 \stackrel{?}{=} y_0$
 $\uparrow \downarrow$ $\downarrow \uparrow$
 $\{a(x_1)\}$ $\{a(y_1), a(y_2)\}$
 $S: k = \{x_1\}$
 $D: (x_1, y_1)$
 $\uparrow \downarrow$ $\downarrow \uparrow$
 $\{b(x_2), c(x_3)\}$ $\{b(y_3)\}$
 $S: k = X$
 $D: \neg _ (_) _ /$

Try with this one



The Seeming Duality

Graded Semantics

D plays (x, y)
Equalities to show
an Equality

S picks a pair $\in R$
 (x', y')

Codensity Liftings

S plays
a Predicate to show
an Inequality

$\hookrightarrow$ D picks a ...

The Seeming Duality

Graded Semantics

D plays
Equalities to show
an Equality

S picks a pair (x', y') $\in R$

(x, y)

Codensity Liftings

S plays
a Predicate to show
an Inequality

~~h~~
~~D picks a ...~~

The Seeming Duality

Graded Semantics

D plays (x, y)
Equalities to show
an Equality

S picks a pair (x', y') $\in R$

Codensity Liftings

S plays $\cup$ $\{e: X \rightarrow Z \text{ are special cases}\}$
Equalities to show
an Inequality

R ~~S~~ D picks a pair (x', y')

The Dual Games

Position	Pl	Move
$(x, y) \in X \times X$	S D	Eq. rel. $R \subseteq X \times X$ s.t. $\gamma(x) \neq_R \gamma(y)$ $\gamma(x) =_R \gamma(y)$
$R \subseteq X \times X$	D S	$(x', y') \notin R$ $(x', y') \in R$

You can alternate non-deterministically!

The Dual Games

Position	Pl	Move
$R \subseteq X \times X$	S D	Eq. rel. $R' \subseteq X \times X$ s.t. $\exists (x, y) \in R$ $\gamma(x) \neq_{R'} \gamma(y)$ $\gamma(x) =_{R'} \gamma(y)$
$R \subseteq X \times X$	D S	$R \subseteq X \times X \setminus R'$ $R \subseteq R'$

The Dual Games

Position	Pl	Move
$R \subseteq X \times X$	$\begin{matrix} S \\ D \end{matrix}$	Eq. rel. $R' \subseteq X \times X$ s.t. $\forall (x, y) \in R \quad \begin{matrix} \gamma(x) \neq_{R'} \gamma(y) \\ \gamma(x) =_{R'} \gamma(y) \end{matrix}$
$R \subseteq X \times X$	$\begin{matrix} D \\ S \end{matrix}$	$\begin{matrix} R \subseteq X \times X \setminus R' \\ R \subseteq R' \end{matrix}$

Singleplayer Game!

TO-DOs

Abstraction?

Relation to Apartness?

Conversion of Winning Strategies?

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