

HAUSDORFF COALGEBRAS

Pedro Nora (collaboration with Dirk Hofmann and Renato Neves)

August 11, 2020

Department of Computer Science, Friedrich-Alexander University
pedro.nora@fau.de

INTRODUCTION

Education (Aveiro)

Education (Aveiro)

- B.Sc. in Mathematics.

Education (Aveiro)

- B.Sc. in Mathematics.
- M.Sc. in Mathematics and Applications (Computer Science).

Education (Aveiro)

- B.Sc. in Mathematics.
- M.Sc. in Mathematics and Applications (Computer Science).
- Ph.D. in Mathematics under the supervision of Dirk Hofmann



Education (Aveiro)

- B.Sc. in Mathematics.
- M.Sc. in Mathematics and Applications (Computer Science).
- Ph.D. in Mathematics under the supervision of Dirk Hofmann
topic: *"Kleisli dualities and Vietoris coalgebras"*

Education (Aveiro)

- B.Sc. in Mathematics.
- M.Sc. in Mathematics and Applications (Computer Science).
- Ph.D. in Mathematics under the supervision of Dirk Hofmann
topic: *"Kleisli dualities and Vietoris coalgebras"*

Previous positions

Education (Aveiro)

- B.Sc. in Mathematics.
- M.Sc. in Mathematics and Applications (Computer Science).
- Ph.D. in Mathematics under the supervision of Dirk Hofmann
topic: *"Kleisli dualities and Vietoris coalgebras"*

Previous positions

- Researcher at INESC TEC

Education (Aveiro)

- B.Sc. in Mathematics.
- M.Sc. in Mathematics and Applications (Computer Science).
- Ph.D. in Mathematics under the supervision of Dirk Hofmann
topic: *"Kleisli dualities and Vietoris coalgebras"*

Previous positions

- Researcher at INESC TEC
- Teacher at IPP

Education (Aveiro)

- B.Sc. in Mathematics.
- M.Sc. in Mathematics and Applications (Computer Science).
- Ph.D. in Mathematics under the supervision of Dirk Hofmann
topic: *"Kleisli dualities and Vietoris coalgebras"*

Previous positions

- Researcher at INESC TEC
- Teacher at IPP

Fun facts

Education (Aveiro)

- B.Sc. in Mathematics.
- M.Sc. in Mathematics and Applications (Computer Science).
- Ph.D. in Mathematics under the supervision of Dirk Hofmann
topic: *"Kleisli dualities and Vietoris coalgebras"*

Previous positions

- Researcher at INESC TEC
- Teacher at IPP

Fun facts

- Alcohol, coffee and coke free

Education (Aveiro)

- B.Sc. in Mathematics.
- M.Sc. in Mathematics and Applications (Computer Science).
- Ph.D. in Mathematics under the supervision of Dirk Hofmann
topic: *"Kleisli dualities and Vietoris coalgebras"*

Previous positions

- Researcher at INESC TEC
- Teacher at IPP

Fun facts

- Alcohol, coffee and coke free
- Linux user

Switching context

- Lutz Schröder and Paul Wild. “Characteristic Logics for Behavioural Metrics via Fuzzy Lax Extensions”. In: *CONCUR* (2020)
- Dirk Hofmann and Pedro Nora. “Hausdorff coalgebras”. In: *Applied Categorical Structures* (2020)

Switching context

- Lutz Schröder and Paul Wild. “Characteristic Logics for Behavioural Metrics via Fuzzy Lax Extensions”. In: *CONCUR* (2020)
- Dirk Hofmann and Pedro Nora. “Hausdorff coalgebras”. In: *Applied Categorical Structures* (2020)

Tholen:

(...) Fuzzy (...)? Perhaps they should work with \mathcal{V} -categories.

QUANTALE-ENRICHED CATEGORY THEORY

WHAT ARE CATEGORIES MADE OF?

Definition

A (small) category \mathbf{C} consists of:

- a set of objects;
- for each pair x,y of objects, a set of morphisms $\mathbf{C}(x, y)$;
- for each triple x,y,z of objects, a function $\circ_{x,y,z}: \mathbf{C}(x, y) \times \mathbf{C}(y, z) \rightarrow \mathbf{C}(x, z)$;
- for each object x , a function $1_x: \mathbf{1} \rightarrow \mathbf{C}(x, x)$;

such that ...

WHAT ARE CATEGORIES MADE OF?

Definition

A (small) category \mathbf{C} consists of:

- a set of objects;
- for each pair x, y of objects, a **set** of morphisms $\mathbf{C}(x, y)$;
- for each triple x, y, z of objects, a **function** $\circ_{x, y, z}: \mathbf{C}(x, y) \times \mathbf{C}(y, z) \rightarrow \mathbf{C}(x, z)$;
- for each object x , a **function** $1_x: \mathbf{1} \rightarrow \mathbf{C}(x, x)$;

such that ...

Definition

A functor $F: \mathbf{X} \rightarrow \mathbf{Y}$ consists of:

- a function from the objects of \mathbf{X} to the objects of \mathbf{Y} ;
- for each pair of objects x, x' of \mathbf{X} a **function** $F_{x, x'}: \mathbf{X}(x, x') \rightarrow \mathbf{Y}(Fx, Fx')$

such that ...

Definition

A **quantale** $\mathcal{V} = (\mathcal{V}, \otimes, k)$ is a complete lattice \mathcal{V} equipped with a commutative monoid structure \otimes , with identity k , such that, for each $u \in \mathcal{V}$,

$$u \otimes - : \mathcal{V} \longrightarrow \mathcal{V} \quad \text{has a right adjoint} \quad \text{hom}(u, -) : \mathcal{V} \longrightarrow \mathcal{V}.$$

Definition

A **quantale** $\mathcal{V} = (\mathcal{V}, \otimes, k)$ is a complete lattice \mathcal{V} equipped with a commutative monoid structure \otimes , with identity k , such that, for each $u \in \mathcal{V}$,

$$u \otimes - : \mathcal{V} \longrightarrow \mathcal{V} \quad \text{has a right adjoint} \quad \text{hom}(u, -) : \mathcal{V} \longrightarrow \mathcal{V}.$$

Recall

- A preordered set (X, \leq) can be understood as a category where for each pair of objects x, y there is at most one morphism between x and y .

A THIN MONOIDAL CATEGORY

Definition

A **quantale** $\mathcal{V} = (\mathcal{V}, \otimes, k)$ is a complete lattice \mathcal{V} equipped with a commutative monoid structure \otimes , with identity k , such that, for each $u \in \mathcal{V}$,

$$u \otimes - : \mathcal{V} \longrightarrow \mathcal{V} \quad \text{has a right adjoint} \quad \text{hom}(u, -) : \mathcal{V} \longrightarrow \mathcal{V}.$$

Recall

- A preordered set (X, \leq) can be understood as a category where for each pair of objects x, y there is at most one morphism between x and y .
- A quantale is a symmetric monoidal closed category.

Category

- a set of objects;

Category

- a set of objects;
- for each pair x, y of objects, a set of morphisms $\mathbf{C}(x, y)$;

Category

- a set of objects;
- for each pair x, y of objects, an **element of \mathcal{V}** , $\mathbf{C}(x, y)$;

DON'T GET LOST IN TRANSLATION

Category

- a set of objects;
- for each pair x, y of objects, an **element of \mathcal{V}** , $\mathbf{C}(x, y)$;
- for each triple x, y, z of objects, a **function** $\circ_{x,y,z}: \mathbf{C}(x, y) \otimes \mathbf{C}(y, z) \rightarrow \mathbf{C}(x, z)$;

Category

- a set of objects;
- for each pair x, y of objects, an **element of \mathcal{V}** , $\mathbf{C}(x, y)$;
- for each triple x, y, z of objects, the **inequality** $\mathbf{C}(x, y) \otimes \mathbf{C}(y, z) \leq \mathbf{C}(x, z)$;

DON'T GET LOST IN TRANSLATION

Category

- a set of objects;
- for each pair x, y of objects, an **element of \mathcal{V}** , $\mathbf{C}(x, y)$;
- for each triple x, y, z of objects, the **inequality** $\mathbf{C}(x, y) \otimes \mathbf{C}(y, z) \leq \mathbf{C}(x, z)$;
- for each object x , a **function** $1_x: \mathbf{1} \rightarrow \mathbf{C}(x, x)$;

DON'T GET LOST IN TRANSLATION

Category

- a set of objects;
- for each pair x, y of objects, an **element of \mathcal{V}** , $\mathbf{C}(x, y)$;
- for each triple x, y, z of objects, the **inequality** $\mathbf{C}(x, y) \otimes \mathbf{C}(y, z) \leq \mathbf{C}(x, z)$;
- for each object x , the **inequality** $k \leq \mathbf{C}(x, x)$;

DON'T GET LOST IN TRANSLATION

Category

- a set of objects;
- for each pair x, y of objects, an **element of \mathcal{V}** , $\mathbf{C}(x, y)$;
- for each triple x, y, z of objects, the **inequality** $\mathbf{C}(x, y) \otimes \mathbf{C}(y, z) \leq \mathbf{C}(x, z)$;
- for each object x , the **inequality** $k \leq \mathbf{C}(x, x)$;

Functor

- a function from the objects of \mathbf{X} to the objects of \mathbf{Y} ;

DON'T GET LOST IN TRANSLATION

Category

- a set of objects;
- for each pair x, y of objects, an **element of \mathcal{V}** , $\mathbf{C}(x, y)$;
- for each triple x, y, z of objects, the **inequality** $\mathbf{C}(x, y) \otimes \mathbf{C}(y, z) \leq \mathbf{C}(x, z)$;
- for each object x , the **inequality** $k \leq \mathbf{C}(x, x)$;

Functor

- a function from the objects of \mathbf{X} to the objects of \mathbf{Y} ;
- for each pair of objects x, x' of \mathbf{X} a **function** $F_{x, x'} : \mathbf{X}(x, x') \rightarrow \mathbf{Y}(Fx, Fx')$

DON'T GET LOST IN TRANSLATION

Category

- a set of objects;
- for each pair x, y of objects, an **element of \mathcal{V}** , $\mathbf{C}(x, y)$;
- for each triple x, y, z of objects, the **inequality** $\mathbf{C}(x, y) \otimes \mathbf{C}(y, z) \leq \mathbf{C}(x, z)$;
- for each object x , the **inequality** $k \leq \mathbf{C}(x, x)$;

Functor

- a function from the objects of \mathbf{X} to the objects of \mathbf{Y} ;
- for each pair of objects x, x' of \mathbf{X} the **inequality** $\mathbf{X}(x, x') \leq \mathbf{Y}(Fx, Fx')$

Definition

Let $\mathcal{V} = (\mathcal{V}, \otimes, k)$ be a quantale.

- A \mathcal{V} -category is a pair (X, a) consisting of a set X and a map $a: X \times X \rightarrow \mathcal{V}$ satisfying

$$k \leq a(x, x) \quad \text{and} \quad a(x, y) \otimes a(y, z) \leq a(x, z),$$

for all $x, y, z \in X$.

- A \mathcal{V} -functor $f: (X, a) \rightarrow (Y, b)$ between \mathcal{V} -categories is a map $f: X \rightarrow Y$ such that

$$a(x, x') \leq b(f(x), f(x')),$$

for all $x, x' \in X$.

- \mathcal{V} -categories and \mathcal{V} -functors define the category $\mathcal{V}\text{-Cat}$.

Example

If \mathcal{V} is the trivial quantale, then $1\text{-Cat} \sim \mathbf{Set}$.

CATEGORIES OF QUANTALE-ENRICHED CATEGORIES

Example

If \mathcal{V} is the trivial quantale, then $1\text{-Cat} \sim \mathbf{Set}$.

Example

If \mathcal{V} is the two element chain $\mathbf{2} = \{0, 1\}$ with $\otimes = \&$. Then $\mathbf{2}\text{-Cat} \sim \mathbf{Ord}$.

CATEGORIES OF QUANTALE-ENRICHED CATEGORIES

Example

If \mathcal{V} is the trivial quantale, then $1\text{-Cat} \sim \mathbf{Set}$.

Example

If \mathcal{V} is the two element chain $\mathbf{2} = \{0, 1\}$ with $\otimes = \&$. Then $\mathbf{2}\text{-Cat} \sim \mathbf{Ord}$.

Example

Consider the quantales based on the extended real half line $\overleftarrow{[0, \infty]}$ ordered by the “greater or equal” relation \geq and

- the tensor product given by addition $+$, denoted by $\overleftarrow{[0, \infty]}_+$;
- or with $\otimes = \max$, denoted as $\overleftarrow{[0, \infty]}_\wedge$.

Then $\overleftarrow{[0, \infty]}_+\text{-Cat} \sim \mathbf{Met}$ is the category of (generalised) metric spaces and non-expansive maps and $\overleftarrow{[0, \infty]}_\wedge\text{-Cat} \sim \mathbf{UMet}$ is the category of (generalised) ultrametric spaces and non-expansive maps.

Example

If \mathcal{V} is the quantale $\overleftarrow{[0, 1]}_{\oplus}$ given by the unit interval $[0, 1]$ with the “greater or equal” relation \geq and the tensor $u \oplus v = \min\{1, u + v\}$. Then $\overleftarrow{[0, 1]}_{\oplus}\text{-Cat} \sim \mathbf{BMet}$ is the category of (generalised) bounded-by-one metric spaces and non-expansive maps.

Example

If \mathcal{V} is the quantale $\overleftarrow{[0, 1]}_{\oplus}$ given by the unit interval $[0, 1]$ with the “greater or equal” relation \geq and the tensor $u \oplus v = \min\{1, u + v\}$. Then $\overleftarrow{[0, 1]}_{\oplus}\text{-Cat} \sim \mathbf{BMet}$ is the category of (generalised) bounded-by-one metric spaces and non-expansive maps.

Alert

Do not forget to reverse the order!

Definition

Let $\mathcal{V} = (\mathcal{V}, \otimes, k)$ be a quantale, and X, Y and Z sets.

- A \mathcal{V} -relation from X to Y , $X \rightsquigarrow Y$, is a map $X \times Y \rightarrow \mathcal{V}$.
- For $r: X \rightsquigarrow Y$ and $s: Y \rightsquigarrow Z$, the **composite** $s \cdot r: X \rightsquigarrow Z$ is calculated pointwise by

$$(s \cdot r)(x, z) = \bigvee_{y \in Y} r(x, y) \otimes s(y, z),$$

for every $x \in X$ and $z \in Z$.

- Sets and \mathcal{V} -relations define the category \mathcal{V} -**Rel**.

Definition

Let $\mathcal{V} = (\mathcal{V}, \otimes, k)$ be a quantale, and X, Y and Z sets.

- A \mathcal{V} -relation from X to Y , $X \rightsquigarrow Y$, is a map $X \times Y \rightarrow \mathcal{V}$.
- For $r: X \rightsquigarrow Y$ and $s: Y \rightsquigarrow Z$, the **composite** $s \cdot r: X \rightsquigarrow Z$ is calculated pointwise by

$$(s \cdot r)(x, z) = \bigvee_{y \in Y} r(x, y) \otimes s(y, z),$$

for every $x \in X$ and $z \in Z$.

- Sets and \mathcal{V} -relations define the category \mathcal{V} -Rel.

Remark

The structure of a \mathcal{V} -category is a reflexive and transitive \mathcal{V} -relation. That is, for a \mathcal{V} -category (X, a) , $1_X \leq a$ and $a \cdot a \leq a$.

COALGEBRAS OF STRICT LIFTINGS

STRICT FUNCTORIAL LIFTINGS

Problem

Given an endofunctor F on a category \mathbf{A} and a faithful functor $U: \mathbf{X} \rightarrow \mathbf{A}$, study a “lifting” of F to an endofunctor \bar{F} on \mathbf{X} . In a strict sense, by “lifting” we mean that the diagram

$$\begin{array}{ccc} \mathbf{X} & \xrightarrow{\bar{F}} & \mathbf{X} \\ \downarrow & & \downarrow \\ \mathbf{A} & \xrightarrow{F} & \mathbf{A} \end{array}$$

commutes.

Problem

Given an endofunctor F on a category \mathbf{A} and a faithful functor $U: \mathbf{X} \rightarrow \mathbf{A}$, study a “lifting” of F to an endofunctor \bar{F} on \mathbf{X} . In a strict sense, by “lifting” we mean that the diagram

$$\begin{array}{ccc} \mathbf{Met} & \xrightarrow{\bar{F}} & \mathbf{Met} \\ \downarrow & & \downarrow \\ \mathbf{Set} & \xrightarrow{F} & \mathbf{Set}. \end{array}$$

commutes^a.

^aPaolo Baldan, Filippo Bonchi, Henning Kerstan, and Barbara König. “Coalgebraic Behavioral Metrics”. In: *Logical Methods in Computer Science* **14**.(3) (2018), pp. 1860–5974.

Problem

Given an endofunctor F on a category \mathbf{A} and a faithful functor $U: \mathbf{X} \rightarrow \mathbf{A}$, study a “lifting” of F to an endofunctor \bar{F} on \mathbf{X} . In a strict sense, by “lifting” we mean that the diagram

$$\begin{array}{ccc} \mathcal{V}\text{-Cat} & \xrightarrow{\bar{F}} & \mathcal{V}\text{-Cat} \\ \downarrow & & \downarrow \\ \mathbf{Set} & \xrightarrow{F} & \mathbf{Set}. \end{array}$$

commutes^a.

^aAdriana Balan, Alexander Kurz, and Jiří Velebil. “Extending set functors to generalised metric spaces”. In: *Logical Methods in Computer Science* **15**.(1) (2019).

HOW TO CONSTRUCT STRICT LIFTINGS?

Lax extension

Consider first a lax extension $\widehat{F}: \mathcal{V}\text{-Rel} \rightarrow \mathcal{V}\text{-Rel}$ of the functor $F: \mathbf{a}$

1. $r \leq r' \implies \widehat{F}r \leq \widehat{F}r'$,
2. $\widehat{F}s \cdot \widehat{F}r \leq \widehat{F}(s \cdot r)$,
3. $Ff \leq \widehat{F}(f)$ and $(Ff)^\circ \leq \widehat{F}(f^\circ)$.

Then, the functor $F: \mathbf{Set} \rightarrow \mathbf{Set}$ admits a natural lifting to $\mathcal{V}\text{-Cat}^b$: the functor $\overline{F}: \mathcal{V}\text{-Cat} \rightarrow \mathcal{V}\text{-Cat}$ sends a \mathcal{V} -category (X, a) to $(FX, \widehat{F}a)$.

^aGavin J. Seal. "Canonical and op-canonical lax algebras". In: *Theory and Applications of Categories* **14**.(10) (2005), pp. 221–243.

^bWalter Tholen. "Ordered topological structures". In: *Topology and its Applications* **156**.(12) (2009), pp. 2148–2157.

Remark

One advantage of this type of lifting is that allows us to use the calculus of \mathcal{V} -relations.

Remark

One advantage of this type of lifting is that allows us to use the calculus of \mathcal{V} -relations.

Proposition

$\bar{F}: \mathcal{V}\text{-Cat} \rightarrow \mathcal{V}\text{-Cat}$ preserves initial \mathcal{V} -functors.

Remark

One advantage of this type of lifting is that allows us to use the calculus of \mathcal{V} -relations.

Proposition

$\bar{F}: \mathcal{V}\text{-Cat} \rightarrow \mathcal{V}\text{-Cat}$ preserves initial \mathcal{V} -functors.

Proof.

Let $f: (X, a) \rightarrow (Y, b)$ be a \mathcal{V} -functor with $a = f^\circ \cdot b \cdot f$. Then $\widehat{F}a = Ff^\circ \cdot \widehat{F}b \cdot Ff$. \square

HOW TO CONSTRUCT LAX EXTENSIONS?

The Wasserstein lifting

One possible way to construct lax extensions based on a (lax) \mathbb{T} -algebra structure $\xi: T\mathcal{V} \rightarrow \mathcal{V}$ is devised in^a: for every \mathcal{V} -relation $r: X \times Y \rightarrow \mathcal{V}$ and for all $x \in TX$ and $\eta \in TY$,

$$\widehat{Tr}(x, \eta) = \bigvee \left\{ \xi \cdot Tr(\mathfrak{w}) \mid \mathfrak{w} \in T(X \times Y), T\pi_1(\mathfrak{w}) = x, T\pi_2(\mathfrak{w}) = \eta \right\}.$$

^aDirk Hofmann. "Topological theories and closed objects". In: *Advances in Mathematics* **215**.(2) (2007), pp. 789–824.

HOW TO CONSTRUCT LAX EXTENSIONS?

The Wasserstein lifting

One possible way to construct lax extensions based on a (lax) \mathbb{T} -algebra structure $\xi: \mathcal{T}\mathcal{V} \rightarrow \mathcal{V}$ is devised in^a: for every \mathcal{V} -relation $r: X \times Y \rightarrow \mathcal{V}$ and for all $x \in TX$ and $\eta \in TY$,

$$\widehat{Tr}(x, \eta) = \bigvee \left\{ \xi \cdot Tr(\mathfrak{w}) \mid \mathfrak{w} \in T(X \times Y), T\pi_1(\mathfrak{w}) = x, T\pi_2(\mathfrak{w}) = \eta \right\}.$$

^aDirk Hofmann. "Topological theories and closed objects". In: *Advances in Mathematics* **215**(2) (2007), pp. 789–824.

Remark

The lax extension \widehat{T} preserves the involution on \mathcal{V} -**Rel**, that is, $\widehat{T}(r^\circ) = (\widehat{Tr})^\circ$ for all \mathcal{V} -relations $r: X \leftrightarrow Y$.

Definition

A functor $U: \mathbf{X} \rightarrow \mathbf{A}$ is called **topological** whenever each cone $\mathcal{C} = (A \rightarrow UX_i)_{i \in I}$ in \mathbf{A} admits a U -initial lifting. In this case we say that \mathbf{X} is topological over \mathbf{A} .

Definition

A functor $U: \mathbf{X} \rightarrow \mathbf{A}$ is called **topological** whenever each cone $\mathcal{C} = (A \rightarrow UX_i)_{i \in I}$ in \mathbf{A} admits a U -initial lifting. In this case we say that \mathbf{X} is topological over \mathbf{A} .

Example

- The category **Top** is topological over **Set**.

Definition

A functor $U: \mathbf{X} \rightarrow \mathbf{A}$ is called **topological** whenever each cone $\mathcal{C} = (A \rightarrow UX_i)_{i \in I}$ in \mathbf{A} admits a U -initial lifting. In this case we say that \mathbf{X} is topological over \mathbf{A} .

Example

- The category **Top** is topological over **Set**.
- For every quantale \mathcal{V} , the categories \mathcal{V} -**Cat** and \mathcal{V} -**Cat**_{sym} are topological over **Set**.

TOPOLOGICAL CATEGORIES

Definition

A functor $U: \mathbf{X} \rightarrow \mathbf{A}$ is called **topological** whenever each cone $\mathcal{C} = (A \rightarrow UX_i)_{i \in I}$ in \mathbf{A} admits a U -initial lifting. In this case we say that \mathbf{X} is topological over \mathbf{A} .

Example

- The category **Top** is topological over **Set**.
- For every quantale \mathcal{V} , the categories \mathcal{V} -**Cat** and \mathcal{V} -**Cat**_{sym} are topological over **Set**.

Proposition

If \mathbf{X} is a topological over a category \mathbf{A} , then \mathbf{X} has limits of shape \mathbf{I} if and only if \mathbf{A} has limits of shape \mathbf{I} .

HOW TO CONSTRUCT STRICT LIFTINGS?

Notation

For a functor $F: \mathbf{A} \rightarrow \mathbf{A}$ and \mathbf{A} -morphisms $\psi: A \rightarrow \tilde{A}$ and $\sigma: F\tilde{A} \rightarrow \tilde{A}$, let $\psi^\diamond: FA \rightarrow \tilde{A}$ denote the composite $FA \xrightarrow{F\psi} F\tilde{A} \xrightarrow{\sigma} \tilde{A}$.

HOW TO CONSTRUCT STRICT LIFTINGS?

Notation

For a functor $F: \mathbf{A} \rightarrow \mathbf{A}$ and \mathbf{A} -morphisms $\psi: A \rightarrow \tilde{A}$ and $\sigma: F\tilde{A} \rightarrow \tilde{A}$, let $\psi^\diamond: FA \rightarrow \tilde{A}$ denote the composite $FA \xrightarrow{F\psi} F\tilde{A} \xrightarrow{\sigma} \tilde{A}$.

The Kantorovich lifting

Consider a category \mathbf{X} equipped with a topological functor $U: \mathbf{X} \rightarrow \mathbf{A}$, and an \mathbf{X} -object \tilde{X} whose underlying set $U\tilde{X}$ carries the structure $\sigma: FUX \rightarrow UX$ of a F -algebra. Then $(\psi^\diamond: FUX \rightarrow UX)_{\psi \in \mathbf{X}(X, \tilde{X})}$ is a U -structured cone, and we define \bar{X} to be the domain of the initial lift of this cone.

Theorem

1. The construction of the previous slide defines a strict lifting $\bar{F}: \mathbf{X} \rightarrow \mathbf{X}$ of the functor $F: \mathbf{A} \rightarrow \mathbf{A}$.
2. For every $\psi: X \rightarrow \tilde{X}$ in \mathbf{X} , ψ^\diamond is an \mathbf{X} -morphism $\psi^\diamond: \bar{F}X \rightarrow \tilde{X}$. In particular, $\sigma = 1_{\tilde{X}}^\diamond$ is an \mathbf{X} -morphism $\sigma: \bar{F}\tilde{X} \rightarrow \tilde{X}$.
3. If \tilde{X} is injective with respect to initial morphisms, then $\bar{F}: \mathbf{X} \rightarrow \mathbf{X}$ preserves initial morphism.
4. Let $\alpha: F \Rightarrow G$ be a natural transformation such that $\sigma_G \cdot \alpha_{\tilde{X}} = \sigma_F$. Then α lifts to a natural transformation between the corresponding \mathbf{X} -functors.
5. If $F = T$ is part of a monad $\mathbb{T} = (T, m, e)$ on \mathbf{A} and $\sigma: T|\tilde{X}| \rightarrow |\tilde{X}|$ is a \mathbb{T} -algebra, then \mathbb{T} lifts naturally to a monad $\bar{\mathbb{T}} = (\bar{T}, m, e)$ on \mathbf{X} .

Theorem

Consider the following commutative diagram of functors.

$$\begin{array}{ccc} \mathbf{X} & \xrightarrow{\bar{F}} & \mathbf{X} \\ \downarrow U & & \downarrow U \\ \mathbf{A} & \xrightarrow{F} & \mathbf{A} \end{array}$$

Theorem

Consider the following commutative diagram of functors.

$$\begin{array}{ccc} \mathbf{X} & \xrightarrow{\bar{F}} & \mathbf{X} \\ \downarrow U & & \downarrow U \\ \mathbf{A} & \xrightarrow{F} & \mathbf{A} \end{array}$$

1. If \bar{F} has a fix-point, then so has F . Hence, if F does not have a fix-point, then neither does \bar{F} .

Theorem

Consider the following commutative diagram of functors.

$$\begin{array}{ccc}
 \mathbf{X} & \xrightarrow{\bar{F}} & \mathbf{X} \\
 \downarrow U & & \downarrow U \\
 \mathbf{A} & \xrightarrow{F} & \mathbf{A}
 \end{array}$$

1. If \bar{F} has a fix-point, then so has F . Hence, if F does not have a fix-point, then neither does \bar{F} .
2. If $U: \mathbf{X} \rightarrow \mathbf{A}$ is topological, then so is $U: \text{CoAlg}(\bar{F}) \rightarrow \text{CoAlg}(F)$.

Theorem

Consider the following commutative diagram of functors.

$$\begin{array}{ccc}
 \mathbf{X} & \xrightarrow{\bar{F}} & \mathbf{X} \\
 \downarrow U & & \downarrow U \\
 \mathbf{A} & \xrightarrow{F} & \mathbf{A}
 \end{array}$$

1. If \bar{F} has a fix-point, then so has F . Hence, if F does not have a fix-point, then neither does \bar{F} .
2. If $U: \mathbf{X} \rightarrow \mathbf{A}$ is topological, then so is $U: \text{CoAlg}(\bar{F}) \rightarrow \text{CoAlg}(F)$.
In particular, the category $\text{CoAlg}(\bar{F})$ has limits of shape I if and only if $\text{CoAlg}(F)$ has limits of shape I .

HAUSDORFF COALGEBRAS

Previous work

Dirk Hofmann, Renato Neves, and Pedro Nora. “Limits in categories of Vietoris coalgebras”. In: *Mathematical Structures in Computer Science* **29**.(4) (2019), pp. 552–587

Previous work

Dirk Hofmann, Renato Neves, and Pedro Nora. “Limits in categories of Vietoris coalgebras”. In: *Mathematical Structures in Computer Science* **29**.(4) (2019), pp. 552–587

Definition

We call a functor **Kripke polynomial** whenever it belongs to the smallest class of endofunctors on **Set** that contains the identity functor, all constant functors and is closed under composition with a powerset functor, products and sums of functors.

Previous work

Dirk Hofmann, Renato Neves, and Pedro Nora. “Limits in categories of Vietoris coalgebras”. In: *Mathematical Structures in Computer Science* **29**.(4) (2019), pp. 552–587

Definition

We call a functor **Kripke polynomial** whenever it belongs to the smallest class of endofunctors on **Set** that contains the identity functor, all constant functors and is closed under composition with a powerset functor, products and sums of functors.

Remark

The class of Kripke polynomial functors is well-behaved in regard to the existence of limits in their respective categories of coalgebras — **if** the powerset functor is submitted to certain cardinality restrictions.

Recall

- The terminal coalgebra for $F: \mathbf{C} \rightarrow \mathbf{C}$ is a fix-point of F .^a

^aJoachim Lambek. “A fixpoint theorem for complete categories”. In: *Mathematische Zeitschrift* **103**.(2) (1968), pp. 151–161.

A BRIEF HISTORY OF “POWERFULL” FUNCTORS

Recall

- The terminal coalgebra for $F: \mathbf{C} \rightarrow \mathbf{C}$ is a fix-point of F .
- The power-set functor $P: \mathbf{Set} \rightarrow \mathbf{Set}$ does not have a fix-point; hence P does not admit a terminal coalgebra.^a

^aGeorg Cantor. “Über eine elementare Frage der Mannigfaltigkeitslehre”. In: *Jahresbericht der Deutschen Mathematiker-Vereinigung* **1** (1891), pp. 75–78.

Recall

- The terminal coalgebra for $F: \mathbf{C} \rightarrow \mathbf{C}$ is a fix-point of F .
- The power-set functor $P: \mathbf{Set} \rightarrow \mathbf{Set}$ does not have a fix-point; hence P does not admit a terminal coalgebra.
- The finite power-set functor $P_{\text{fin}}: \mathbf{Set} \rightarrow \mathbf{Set}$ admits a terminal coalgebra (for instance, because P_{fin} is finitary).^a

^aMichael Barr. “Terminal coalgebras in well-founded set theory”. In: *Theoretical Computer Science* 114.(2) (1993), pp. 299–315.

A BRIEF HISTORY OF “POWERFULL” FUNCTORS

Recall

- The terminal coalgebra for $F: \mathbf{C} \rightarrow \mathbf{C}$ is a fix-point of F .
- The power-set functor $P: \mathbf{Set} \rightarrow \mathbf{Set}$ does not have a fix-point; hence P does not admit a terminal coalgebra.
- The finite power-set functor $P_{\text{fin}}: \mathbf{Set} \rightarrow \mathbf{Set}$ admits a terminal coalgebra (for instance, because P_{fin} is finitary).
- *Somehow more general*: the Vietoris functor $V: \mathbf{CompHaus} \rightarrow \mathbf{CompHaus}$ admits a terminal coalgebra^a

^aHere: V preserves codirected limits. This result appears as an exercise in Ryszard Engelking. *General topology*. 2nd ed. Vol. 6. Sigma Series in Pure Mathematics. Berlin: Heldermann Verlag, 1989. viii + 529.

A BRIEF HISTORY OF “POWERFULL” FUNCTORS

Recall

- The terminal coalgebra for $F: \mathbf{C} \rightarrow \mathbf{C}$ is a fix-point of F .
- The power-set functor $P: \mathbf{Set} \rightarrow \mathbf{Set}$ does not have a fix-point; hence P does not admit a terminal coalgebra.
- The finite power-set functor $P_{\text{fin}}: \mathbf{Set} \rightarrow \mathbf{Set}$ admits a terminal coalgebra (for instance, because P_{fin} is finitary).
- *Somehow more general*: the Vietoris functor $V: \mathbf{CompHaus} \rightarrow \mathbf{CompHaus}$ admits a terminal coalgebra (and the same is true for $V: \mathbf{PosComp} \rightarrow \mathbf{PosComp}$).^{ab}

^aLeopoldo Nachbin. *Topologia e Ordem*. University of Chicago Press, 1950.

^bDirk Hofmann, Renato Neves, and Pedro Nora. “Limits in categories of Vietoris coalgebras”. In: *Mathematical Structures in Computer Science* 20 (1) (2010), pp. F52–F87.

A BRIEF HISTORY OF “POWERFULL” FUNCTORS

Recall

- The terminal coalgebra for $F: \mathbf{C} \rightarrow \mathbf{C}$ is a fix-point of F .
- The power-set functor $P: \mathbf{Set} \rightarrow \mathbf{Set}$ does not have a fix-point; hence P does not admit a terminal coalgebra.
- The finite power-set functor $P_{\text{fin}}: \mathbf{Set} \rightarrow \mathbf{Set}$ admits a terminal coalgebra (for instance, because P_{fin} is finitary).
- *Somehow more general:* the Vietoris functor $V: \mathbf{CompHaus} \rightarrow \mathbf{CompHaus}$ admits a terminal coalgebra (and the same is true for $V: \mathbf{PosComp} \rightarrow \mathbf{PosComp}$).
- *A bit more general:* the compact Vietoris functor $V_c: \mathbf{Top} \rightarrow \mathbf{Top}$ admits a terminal coalgebra.

A BRIEF HISTORY OF “POWERFULL” FUNCTORS

Recall

- The terminal coalgebra for $F: \mathbf{C} \rightarrow \mathbf{C}$ is a fix-point of F .
- The power-set functor $P: \mathbf{Set} \rightarrow \mathbf{Set}$ does not have a fix-point; hence P does not admit a terminal coalgebra.
- The finite power-set functor $P_{\text{fin}}: \mathbf{Set} \rightarrow \mathbf{Set}$ admits a terminal coalgebra (for instance, because P_{fin} is finitary).
- *Somehow more general:* the Vietoris functor $V: \mathbf{CompHaus} \rightarrow \mathbf{CompHaus}$ admits a terminal coalgebra (and the same is true for $V: \mathbf{PosComp} \rightarrow \mathbf{PosComp}$).
- *A bit more general:* the compact Vietoris functor $V_c: \mathbf{Top} \rightarrow \mathbf{Top}$ admits a terminal coalgebra.
- *A bit surprising(?):* Also the lower Vietoris functor $V: \mathbf{Top} \rightarrow \mathbf{Top}$ admits a terminal coalgebra.

Definition

Let $f: (X, a) \rightarrow (Y, b)$ be a \mathcal{V} -functor.

1. For every $A \subseteq X$, put $\uparrow^a A = \{y \in X \mid k \leq \bigvee_{x \in A} a(x, y)\}$.

Definition

Let $f: (X, a) \rightarrow (Y, b)$ be a \mathcal{V} -functor.

1. For every $A \subseteq X$, put $\uparrow^a A = \{y \in X \mid k \leq \bigvee_{x \in A} a(x, y)\}$.
2. We call a subset $A \subseteq X$ of (X, a) **increasing** whenever $A = \uparrow^a A$.

Definition

Let $f: (X, a) \rightarrow (Y, b)$ be a \mathcal{V} -functor.

1. For every $A \subseteq X$, put $\uparrow^a A = \{y \in X \mid k \leq \bigvee_{x \in A} a(x, y)\}$.
2. We call a subset $A \subseteq X$ of (X, a) **increasing** whenever $A = \uparrow^a A$.
3. We consider the \mathcal{V} -category $HX = \{A \subseteq X \mid A \text{ is increasing}\}$, equipped with $\text{Ha}(A, B) = \bigwedge_{y \in B} \bigvee_{x \in A} a(x, y)$, for all $A, B \in HX$.

Definition

Let $f: (X, a) \rightarrow (Y, b)$ be a \mathcal{V} -functor.

1. For every $A \subseteq X$, put $\uparrow^a A = \{y \in X \mid k \leq \bigvee_{x \in A} a(x, y)\}$.
2. We call a subset $A \subseteq X$ of (X, a) **increasing** whenever $A = \uparrow^a A$.
3. We consider the \mathcal{V} -category $HX = \{A \subseteq X \mid A \text{ is increasing}\}$, equipped with $Ha(A, B) = \bigwedge_{y \in B} \bigvee_{x \in A} a(x, y)$, for all $A, B \in HX$.
4. The map $Hf: H(X, a) \rightarrow H(Y, b)$ sends an increasing subset $A \subseteq X$ to $\uparrow^b f(A)$.

Definition

Let $f: (X, a) \rightarrow (Y, b)$ be a \mathcal{V} -functor.

1. For every $A \subseteq X$, put $\uparrow^a A = \{y \in X \mid k \leq \bigvee_{x \in A} a(x, y)\}$.
2. We call a subset $A \subseteq X$ of (X, a) **increasing** whenever $A = \uparrow^a A$.
3. We consider the \mathcal{V} -category $HX = \{A \subseteq X \mid A \text{ is increasing}\}$, equipped with $Ha(A, B) = \bigwedge_{y \in B} \bigvee_{x \in A} a(x, y)$, for all $A, B \in HX$.
4. The map $Hf: H(X, a) \rightarrow H(Y, b)$ sends an increasing subset $A \subseteq X$ to $\uparrow^b f(A)$.
5. The functor H is part of a Kock–Zöberlein monad $\mathbb{H} = (H, w, \hat{h})$ on \mathcal{V} -Cat.

$$\hat{h}_X: X \longrightarrow HX,$$

$$x \longmapsto \uparrow x$$

$$w_X: HHX \longrightarrow HX.$$

$$A \longmapsto \bigcup A$$

Theorem

Let X be a partially ordered set. Then there is no embedding $\varphi: \text{Up}(X) \rightarrow X$.^a

^aRobert P. Dilworth and Andrew M. Gleason. "A generalized Cantor theorem". In: *Proceedings of the American Mathematical Society* **13**.(5) (1962), pp. 704–705.

Theorem

Let X be a partially ordered set. Then there is no embedding $\varphi: \text{Up}(X) \rightarrow X$.^a

^aRobert P. Dilworth and Andrew M. Gleason. "A generalized Cantor theorem". In: *Proceedings of the American Mathematical Society* **13**.(5) (1962), pp. 704–705.

Corollary

The upset functor $\text{Up}: \mathbf{Ord} \rightarrow \mathbf{Ord}$ does not admit a terminal coalgebra.

Theorem

Let X be a partially ordered set. Then there is no embedding $\varphi: \text{Up}(X) \rightarrow X$.^a

^aRobert P. Dilworth and Andrew M. Gleason. "A generalized Cantor theorem". In: *Proceedings of the American Mathematical Society* **13**.(5) (1962), pp. 704–705.

Corollary

The upset functor $\text{Up}: \mathbf{Ord} \rightarrow \mathbf{Ord}$ does not admit a terminal coalgebra.

Remark

The category $\text{CoAlg}(\text{Up})$ has equalisers.

Theorem

Let \mathcal{V} be a non-trivial quantale and (X, a) be a \mathcal{V} -category. There is no embedding of type $H(X, a) \rightarrow (X, a)$.

Theorem

Let \mathcal{V} be a non-trivial quantale and (X, a) be a \mathcal{V} -category. There is no embedding of type $H(X, a) \rightarrow (X, a)$.

Corollary

Let \mathcal{V} be a non-trivial quantale. The Hausdorff functor $H: \mathcal{V}\text{-Cat} \rightarrow \mathcal{V}\text{-Cat}$ does not admit a terminal coalgebra, neither does any possible restriction to a full subcategory of $\mathcal{V}\text{-Cat}$.

Theorem

Let \mathcal{V} be a non-trivial quantale and (X, a) be a \mathcal{V} -category. There is no embedding of type $H(X, a) \rightarrow (X, a)$.

Corollary

Let \mathcal{V} be a non-trivial quantale. The Hausdorff functor $H: \mathcal{V}\text{-Cat} \rightarrow \mathcal{V}\text{-Cat}$ does not admit a terminal coalgebra, neither does any possible restriction to a full subcategory of $\mathcal{V}\text{-Cat}$.

Remark

In particular, the (non-symmetric) Hausdorff functor on **Met** does not admit a terminal coalgebra, and the same applies to its restriction to the full subcategory of compact metric spaces.

Theorem

Let \mathcal{V} be a non-trivial quantale and (X, a) be a \mathcal{V} -category. There is no embedding of type $H(X, a) \rightarrow (X, a)$.

Corollary

Let \mathcal{V} be a non-trivial quantale. The Hausdorff functor $H: \mathcal{V}\text{-Cat} \rightarrow \mathcal{V}\text{-Cat}$ does not admit a terminal coalgebra, neither does any possible restriction to a full subcategory of $\mathcal{V}\text{-Cat}$.

Remark

In particular, the (non-symmetric) Hausdorff functor on **Met** does not admit a terminal coalgebra, and the same applies to its restriction to the full subcategory of compact metric spaces. Passing to the symmetric version of the Hausdorff functor does not remedy the situation.

ADDING TOPOLOGY

Extending the Ultrafilter monad

We assume that \mathcal{V} is a completely distributive quantale, then

$$\xi: \mathbb{U}\mathcal{V} \longrightarrow \mathcal{V}, \quad \mathfrak{v} \longmapsto \bigwedge_{A \in \mathfrak{v}} \bigvee A$$

is the structure of an \mathbb{U} -algebra on \mathcal{V} (the Lawson topology).

Extending the Ultrafilter monad

We assume that \mathcal{V} is a completely distributive quantale, then

$$\xi: \mathbb{U}\mathcal{V} \longrightarrow \mathcal{V}, \quad \mathfrak{v} \longmapsto \bigwedge_{A \in \mathfrak{v}} \bigvee A$$

is the structure of an \mathbb{U} -algebra on \mathcal{V} (the Lawson topology). Therefore we obtain a lax extension of the ultrafilter monad to \mathcal{V} -**Rel** that induces a monad on \mathcal{V} -**Cat**.

Here:

$$\mathbb{U}a(x, \eta) = \bigwedge_{A, B} \bigvee_{x, y} a(x, y), \quad (X, a) \longmapsto (\mathbb{U}X, \mathbb{U}a).$$

Extending the Ultrafilter monad

We assume that \mathcal{V} is a completely distributive quantale, then

$$\xi: \mathbb{U}\mathcal{V} \longrightarrow \mathcal{V}, \quad v \longmapsto \bigwedge_{A \in v} \bigvee A$$

is the structure of an \mathbb{U} -algebra on \mathcal{V} (the Lawson topology). Therefore we obtain a lax extension of the ultrafilter monad to $\mathcal{V}\text{-Rel}$ that induces a monad on $\mathcal{V}\text{-Cat}$.

Its algebras are \mathcal{V} -categories equipped with a *compatible* compact Hausdorff topology^{ab}; we call them **\mathcal{V} -categorical compact Hausdorff spaces**, and denote the corresponding Eilenberg–Moore category by $\mathcal{V}\text{-CatCH}$.

^aLeopoldo Nachbin. *Topologia e Ordem*. University of Chicago Press, 1950.

^bWalter Tholen. “Ordered topological structures”. In: *Topology and its Applications* **156**.(12) (2009), pp. 2148–2157.

Extending the Ultrafilter monad

We assume that \mathcal{V} is a completely distributive quantale, then

$$\xi: \mathbb{U}\mathcal{V} \longrightarrow \mathcal{V}, \quad v \longmapsto \bigwedge_{A \in v} \bigvee A$$

is the structure of an \mathbb{U} -algebra on \mathcal{V} (the Lawson topology). Therefore we obtain a lax extension of the ultrafilter monad to $\mathcal{V}\text{-Rel}$ that induces a monad on $\mathcal{V}\text{-Cat}$.

Its algebras are \mathcal{V} -categories equipped with a *compatible* compact Hausdorff topology; we call them **\mathcal{V} -categorical compact Hausdorff spaces**, and denote the corresponding Eilenberg–Moore category by $\mathcal{V}\text{-CatCH}$.

Theorem

For an ordered set (X, \leq) and a \mathbb{U} -algebra (X, α) , the following are equivalent.

- (i) $\alpha: (\mathbb{U}X, \mathbb{U}\leq) \rightarrow (X, \leq)$ is monotone.

Extending the Ultrafilter monad

We assume that \mathcal{V} is a completely distributive quantale, then

$$\xi: \mathbb{U}\mathcal{V} \longrightarrow \mathcal{V}, \quad v \longmapsto \bigwedge_{A \in v} \bigvee A$$

is the structure of an \mathbb{U} -algebra on \mathcal{V} (the Lawson topology). Therefore we obtain a lax extension of the ultrafilter monad to $\mathcal{V}\text{-Rel}$ that induces a monad on $\mathcal{V}\text{-Cat}$.

Its algebras are \mathcal{V} -categories equipped with a *compatible* compact Hausdorff topology; we call them **\mathcal{V} -categorical compact Hausdorff spaces**, and denote the corresponding Eilenberg–Moore category by $\mathcal{V}\text{-CatCH}$.

Theorem

For an ordered set (X, \leq) and a \mathbb{U} -algebra (X, α) , the following are equivalent.

- (i) $\alpha: (\mathbb{U}X, \mathbb{U}\leq) \rightarrow (X, \leq)$ is monotone.
- (ii) $G_{\leq} \subseteq X \times X$ is closed.

Extending the Ultrafilter monad

We assume that \mathcal{V} is a completely distributive quantale, then

$$\xi: \mathbb{U}\mathcal{V} \longrightarrow \mathcal{V}, \quad v \longmapsto \bigwedge_{A \in v} \bigvee A$$

is the structure of an \mathbb{U} -algebra on \mathcal{V} (the Lawson topology). Therefore we obtain a lax extension of the ultrafilter monad to $\mathcal{V}\text{-Rel}$ that induces a monad on $\mathcal{V}\text{-Cat}$.

Its algebras are \mathcal{V} -categories equipped with a *compatible* compact Hausdorff topology; we call them **\mathcal{V} -categorical compact Hausdorff spaces**, and denote the corresponding Eilenberg–Moore category by $\mathcal{V}\text{-CatCH}$.

Theorem

For a \mathcal{V} -category (X, a) and a \mathbb{U} -algebra (X, α) , the following are equivalent.

- (i) $\alpha: \mathbb{U}(X, a) \rightarrow (X, a)$ is a \mathcal{V} -functor.
- (ii) $a: (X, \alpha) \times (X, \alpha) \rightarrow (\mathcal{V}, \xi_{\leq})$ is continuous.

Lemma

Let (X, a, α) be a \mathcal{V} -categorical compact Hausdorff space and $A, B \subseteq X$ so that $A \cap B = \emptyset$, A is increasing and compact in $(X, \alpha_{\leq})^{\text{op}}$ and B is compact in (X, α_{\leq}) . Then there exists some $u \ll k$ so that, for all $x \in A$ and $y \in B$, $u \not\leq a(x, y)$.

Lemma

Let (X, a, α) be a \mathcal{V} -categorical compact Hausdorff space and $A, B \subseteq X$ so that $A \cap B = \emptyset$, A is increasing and compact in $(X, \alpha_{\leq})^{\text{op}}$ and B is compact in (X, α_{\leq}) . Then there exists some $u \ll k$ so that, for all $x \in A$ and $y \in B$, $u \not\leq a(x, y)$.

Corollary

For every compact subset $A \subseteq X$ of $(X, \alpha_{\leq})^{\text{op}}$, $\uparrow^a A = \uparrow^{\leq} A$. In particular, for every closed subset $A \subseteq X$ of (X, α) , $\uparrow^a A = \uparrow^{\leq} A$.

Lemma

Let (X, a, α) be a \mathcal{V} -categorical compact Hausdorff space and $A, B \subseteq X$ so that $A \cap B = \emptyset$, A is increasing and compact in $(X, \alpha_{\leq})^{\text{op}}$ and B is compact in (X, α_{\leq}) . Then there exists some $u \ll k$ so that, for all $x \in A$ and $y \in B$, $u \not\leq a(x, y)$.

Corollary

For every compact subset $A \subseteq X$ of $(X, \alpha_{\leq})^{\text{op}}$, $\uparrow^a A = \uparrow^{\leq} A$. In particular, for every closed subset $A \subseteq X$ of (X, α) , $\uparrow^a A = \uparrow^{\leq} A$.

Theorem (Nachbin)

Let $A \subseteq X$ be closed and decreasing and $B \subseteq X$ be closed and increasing with $A \cap B = \emptyset$. Then there exist $V \subseteq X$ open and co-increasing and $W \subseteq X$ open and co-decreasing with

$$A \subseteq V, \quad B \subseteq W, \quad V \cap W = \emptyset.$$

THE HAUSDORFF MONAD (AGAIN)

Definition

For a \mathcal{V} -categorical compact Hausdorff space $X = (X, a, \alpha)$, we put

$$HX = \{A \subseteq X \mid A \text{ is closed and increasing}\}$$

with the restriction of the Hausdorff structure to HX and the **hit-and-miss topology** (Vietoris topology). That is, the topology generated by the sets

$$V^\diamond = \{A \in HX \mid A \cap V \neq \emptyset\} \quad (V \text{ open, co-increasing})$$

and

$$W^\square = \{A \in HX \mid A \subseteq W\} \quad (W \text{ open, co-decreasing}).$$

THE HAUSDORFF MONAD (AGAIN)

Definition

For a \mathcal{V} -categorical compact Hausdorff space $X = (X, a, \alpha)$, we put

$$HX = \{A \subseteq X \mid A \text{ is closed and increasing}\}$$

with the restriction of the Hausdorff structure to HX and the **hit-and-miss topology** (Vietoris topology).

Proposition

For every \mathcal{V} -categorical compact Hausdorff space X , HX is a \mathcal{V} -categorical compact Hausdorff space.

Compare with:

For a compact metric space, the Hausdorff metric induces the Vietoris topology.

THE HAUSDORFF MONAD (AGAIN)

Definition

For a \mathcal{V} -categorical compact Hausdorff space $X = (X, a, \alpha)$, we put

$$HX = \{A \subseteq X \mid A \text{ is closed and increasing}\}$$

with the restriction of the Hausdorff structure to HX and the **hit-and-miss topology** (Vietoris topology).

Proposition

For every \mathcal{V} -categorical compact Hausdorff space X , HX is a \mathcal{V} -categorical compact Hausdorff space.

Theorem

The construction above defines a functor $H: \mathcal{V}\text{-CatCH} \rightarrow \mathcal{V}\text{-CatCH}$.

THE HAUSDORFF MONAD (AGAIN)

Definition

For a \mathcal{V} -categorical compact Hausdorff space $X = (X, a, \alpha)$, we put

$$HX = \{A \subseteq X \mid A \text{ is closed and increasing}\}$$

with the restriction of the Hausdorff structure to HX and the **hit-and-miss topology** (Vietoris topology).

Proposition

For every \mathcal{V} -categorical compact Hausdorff space X , HX is a \mathcal{V} -categorical compact Hausdorff space.

Theorem

The construction above defines a functor $H: \mathcal{V}\text{-CatCH} \rightarrow \mathcal{V}\text{-CatCH}$.

In fact, we obtain a Kock-Zöberlein monad.

Proposition

The diagrams of functors commutes.

$$\begin{array}{ccc}
 \mathbf{OrdCH} & \xrightarrow{H} & \mathbf{OrdCH} \\
 \left(\begin{array}{c} \downarrow \\ \uparrow \end{array} \right) & & \left(\begin{array}{c} \downarrow \\ \uparrow \end{array} \right) \\
 \mathcal{V}\text{-CatCH} & \xrightarrow{H} & \mathcal{V}\text{-CatCH}
 \end{array}$$

Proposition

The diagrams of functors commutes.

$$\begin{array}{ccc}
 \mathbf{OrdCH} & \xrightarrow{H} & \mathbf{OrdCH} \\
 \left(\begin{array}{c} \uparrow \\ \downarrow \end{array} \right) & & \left(\begin{array}{c} \uparrow \\ \downarrow \end{array} \right) \\
 \mathcal{V}\text{-CatCH} & \xrightarrow{H} & \mathcal{V}\text{-CatCH}
 \end{array}$$

Proposition

The Hausdorff functor on $\mathcal{V}\text{-CatCH}$ preserves codirected initial cones with respect to the forgetful functor $\mathcal{V}\text{-CatCH} \rightarrow \mathbf{CompHaus}$.

Proposition

The diagrams of functors commutes.

$$\begin{array}{ccc}
 \mathbf{OrdCH} & \xrightarrow{H} & \mathbf{OrdCH} \\
 \left(\begin{array}{c} \uparrow \\ \downarrow \end{array} \right) & & \left(\begin{array}{c} \uparrow \\ \downarrow \end{array} \right) \\
 \mathcal{V}\text{-CatCH} & \xrightarrow{H} & \mathcal{V}\text{-CatCH}
 \end{array}$$

Proposition

The Hausdorff functor on $\mathcal{V}\text{-CatCH}$ preserves codirected initial cones with respect to the forgetful functor $\mathcal{V}\text{-CatCH} \rightarrow \mathbf{CompHaus}$.

Theorem

The Hausdorff functor $H: \mathcal{V}\text{-CatCH} \rightarrow \mathcal{V}\text{-CatCH}$ preserves codirected limits.

Theorem

For $H: \mathcal{V}\text{-CatCH} \rightarrow \mathcal{V}\text{-CatCH}$, the forgetful functor $\text{CoAlg}(H) \rightarrow \mathcal{V}\text{-CatCH}$ is comonadic. Moreover, $\mathcal{V}\text{-CatCH}$ has equalisers and is therefore complete.

COALGEBRAS FOR THE HAUSDORFF FUNCTOR

Theorem

For $H: \mathcal{V}\text{-Cat}\mathbf{CH} \rightarrow \mathcal{V}\text{-Cat}\mathbf{CH}$, the forgetful functor $\text{CoAlg}(H) \rightarrow \mathcal{V}\text{-Cat}\mathbf{CH}$ is comonadic. Moreover, $\mathcal{V}\text{-Cat}\mathbf{CH}$ has equalisers and is therefore complete.

Theorem

The category of coalgebras of a Hausdorff polynomial functor on $\mathcal{V}\text{-Cat}\mathbf{CH}$ is (co)complete.

Definition

We call a functor **Hausdorff polynomial** whenever it belongs to the smallest class of endofunctors on $\mathcal{V}\text{-Cat}$ that contains the identity functor, all constant functors and is closed under composition with H , products and sums of functors.